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The quantum comb:

theory and applications

to quantum networks

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Theory of Quantum Combs in collaboration with



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Applications to Optimal

Q-Tomography and Q-Learning



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New Quantum Estimation Theory, with multiple copies, and optimization of the setup





Outline

- New Quantum Estimation Theory, with multiple copies, and optimization of the setup
 - Convex optimization method based on the new notions of quantum comb and quantum tester
- Applications:
 - discrimination of unitary operators and of memory channels (quantum oracle-calling algoritms)



- strategies in quantum protocols, crypto and games
- quantum-algorithm learning (storing undisclosable unitaries)
 - process-cloning
 - optimal tomography



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Quantum Estimation Theory

Quantum state $ho_{ heta}$ parameterized by heta

Problem: estimate θ optimally according to the cost function $C(\theta, \hat{\theta})$





Quantum Estimation Theory

Quantum state $ho_{ heta}$ parameterized by heta

Problem: estimate θ optimally according to the cost function $C(\theta, \hat{\theta})$

Mathematical formulation:

find the optimal POVM $P_{\hat{\theta}}$ minimizing the cost





Quantum Estimation Theory

Practically interesting situation (e.g. for the phase of an e.m. mode):

 $\theta \Longrightarrow \rho_{\theta} = U_{\theta} \rho U_{\theta}^{\dagger}$





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Then you want also to optimize ho

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Quantum Estimation Theory

Practically interesting situation (e.g. for the phase of an e.m. mode):

Then you want also to optimize ρ

The optimal POVM for estimating $\theta\,$ depends on ρ

 $\theta \Longrightarrow \rho_{\theta} = U_{\theta} \rho U_{\theta}^{\dagger}$





Quantum Estimation Theory

 $= P_{\hat{\theta}}$

Practically interesting situation (e.g. for the phase of an e.m. mode):

Then you want also to optimize ho

The optimal POVM for estimating θ depends on ρ

 $\theta \Longrightarrow \rho_{\theta} = U_{\theta} \rho U_{\theta}^{\dagger}$

Interesting situation: the parameter to be estimated is encoded on a transformation---not on the state!



Problem: estimate x parameterizing the (unitary) transformation U_x optimally according to the cost function



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Lesson that we learned from entanglement:





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Find the optimal entangled state $|\Psi\rangle\rangle$ (with an any possible ancilla) along with the optimal joint POVM $P_{\mathcal{X}}$





Problem: estimate x parameterizing the (unitary) transformation U_x optimally according to the cost function

Lesson that we learned from entanglement:

Find the optimal entangled state $|\Psi\rangle\rangle$ (with an any possible ancilla) along with the optimal joint POVM P_x

For the phase no need of entanglement (we were lucky!)







New scheme









quantum feedback: perform a transformation \mathscr{T}_U on a system which depends on an unknown unitary transformation U occurring on a (generally different) system (e.g. reference-frames realignment).







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Multiple copies

 \clubsuit For parameter estimation: repeat the estimation N times, gaining a precision factor \sqrt{N}



Multiple copies

- \checkmark For parameter estimation: repeat the estimation N times, gaining a precision factor \sqrt{N}
- We have better use a coherent strategy, in which you perform a joint POVM



Multiple copies

- \checkmark For parameter estimation: repeat the estimation N times, gaining a precision factor \sqrt{N}
- Weight However, you better use a coherent strategy, in which you perform a joint POVM

and you want to do the same for the quantum feedback



What is the best that you can do?

Use a Quantum Board!



General scheme: put the copies of the unknown unitary in a suitable quantum circuit which performs the desired transformation/estimation.

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Quantum circuit board: input and output are themselves circuits that are slotted into the board.



Use a Quantum Board!



Optimize the quantum circuit board for all possible dispositions of the slots



It looks a difficult problem ...



In parallel over a joint entangled state?





In sequence intercalated by some unitary?



In parallel over a joint entangled state?





In sequence intercalated by some unitary?



In parallel over a joint entangled state?



Asymptotically: same sensitivity [Giovannetti, LLoyd, Maccone, PRL 96, 010401 (2006)]

$-\underbrace{U_{\phi}}_{\phi} - \underbrace{U_{\phi}}_{\phi} - \underbrace{U$

In sequence intercalated by some unitary?

For unitary discrimination:[Duan, Feng, Ying, PRL 98, 100503 (2007)]

In parallel over a joint entangled state?

For unitary discrimination: G.M.D'Ariano, P. Lo Presti, M. Paris, PRL 87, 270404 (2001); A. Acín, E. Jané, and G. Vidal, Phys. Rev. A 64, 050302 (2001)

U_φ

Asymptotically: same sensitivity [Giovannetti, LLoyd, Maccone, PRL 96, 010401 (2006)]

An optimal board architecture [van Dam, D'Ariano, Ekert, Macchiavello, Mosca, PRL 98, 090501 (2007)]


For example: what is the optimal board for phase estimation?

An optimal board architecture [van Dam, D'Ariano, Ekert, Macchiavello, Mosca, PRL 98, 090501 (2007)]





What is the mathematical formulation of the problem?





It can be regarded as an equivalence class of quantum circuits performing the same input-output transformation ...







It can be regarded as an equivalence class of quantum circuits performing the same input-output transformation ... For a channel the input and the output are states







Equivalence class of quantum circuits boards performing the same overall input-output transformation ...







Equivalence class of quantum circuits boards performing the same overall input-output transformation ...

But now, the input and the output are transformations





Quantum Board



Problem: what is the optimal board for given slots achieving a global input/output transformation optimally according to a given cost function?



Quantum Combs

G.Chiribella, G.M.D'Ariano, P.Perinotti, PRL 101 060401 (2008)

All circuits-boards can be reshaped in form of "combs", with an ordered sequence of slots, each between two successive teeth





Quantum Combs

PRL 101 060401 (2008)

G.Chiribella, G.M.D'Ariano, P.Perinotti, PRL 101 060401 (2008)



Pins = quantum systems with generally variable dimensions



How do we describe a quantum comb mathematically?

Channel: Choi representation

Mathematically the input-output transformation operated by a quantum circuit is a CP map, and is in one-to-one correspondence with a positive operator called "Choi-Jamiolkowski operator"---the output state of the map applied locally to a maximally entangled state.





Causal networks

The quantum comb is equivalent to a causal network with all inputs on the left and all outputs on the right



QUIT quantum informatio theory grou PRL 101 060401 (2008)

Causal networks

The quantum comb is equivalent to a causal network with all inputs on the left and all outputs on the right



The causal network is also equivalent to the stack of memory channels







PRL 101 060401 (2008) Choi representation





Choi representation

PRL 101 060401 (2008)



Causality constraints: (N+1 inputs/outputs)

$$\operatorname{Tr}_{2n-1}[R^{(n)}] = I_{2n-2} \otimes R^{(n-1)}, \quad n = 1, \dots, N$$

 $R^{(0)} = 1, \quad R^{(n)} = R$





A quantum comb performs a transformation that is a generalization of the quantum operation: the so called "supermap"



A supermap sends a series of N channels to one channel, also when applied locally, e.g.





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Mathematically it is represented by a CP N-linear map which sends N Choi operators to one Choi operator, and with his own Choi operator satisfying the causality constraints.





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(we can likewise consider probabilistic supermaps).





More generally, a quantum comb maps a series of channels into a comb







More generally, a quantum comb maps a series of channels into a comb



or, even more generally, a comb to a comb







The notion of supermap is the last level of generalization, i.e. "super-supermaps" (mapping supermaps to supermaps) are still supermaps = quantum combs.





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$A_{in} \begin{bmatrix} a & f & c \\ b & d & d \end{bmatrix} A_{out} \implies H_{in} \begin{bmatrix} a & f & c \\ b & d & f \\ e & g \end{bmatrix} H_{out}$ $B_{in} \begin{bmatrix} d & f & f \\ e & g & g \end{bmatrix} B_{out} \qquad Choi-operator calculus$ $A \in B(A_{out} \otimes A_{in}) = B(H_a \otimes H_b \otimes H_c \otimes H_d), \qquad J \equiv H_d$ $B \in B(B_{out} \otimes B_{in}) = B(H_d \otimes H_e \otimes H_f \otimes H_g)$



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$\mathsf{A}_{\mathrm{in}} \begin{bmatrix} a \\ b \end{bmatrix} \mathscr{A} \begin{bmatrix} -c \\ -d \end{bmatrix} \mathsf{A}_{\mathrm{out}}$ $\begin{array}{c|c} \mathbf{H}_{\mathrm{in}} & a & c \\ b & \mathbf{\mathcal{A}} & \mathbf{\mathcal{A}} & f \\ e & \mathbf{\mathcal{B}} & g \end{array} \end{array} \begin{array}{c} \mathbf{H}_{\mathrm{out}} \\ \mathbf{H}_{\mathrm{out}} \end{array}$ $\mathsf{B}_{\mathrm{in}} \begin{bmatrix} \mathbf{d} & & \\ e & & \mathcal{B} \end{bmatrix} = \begin{bmatrix} f \\ g \end{bmatrix} \mathsf{B}_{\mathrm{out}}$ Choi-operator calculus $A \in \mathsf{B}(\mathsf{A}_{\mathrm{out}} \otimes \mathsf{A}_{\mathrm{in}}) = \mathsf{B}(\mathsf{H}_a \otimes \mathsf{H}_b \otimes \mathsf{H}_c \otimes \mathsf{H}_d),$ $J \equiv H_d$ $B \in \mathsf{B}(\mathsf{B}_{\mathrm{out}} \otimes \mathsf{B}_{\mathrm{in}}) = \mathsf{B}(\mathsf{H}_{d} \otimes \mathsf{H}_{e} \otimes \mathsf{H}_{f} \otimes \mathsf{H}_{q})$ $AB := (A \otimes I_{e,f,q})(I_{a,b,c} \otimes B)$



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PRL 101 060401 (2008)

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The link-product is commutative!



Link product





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Special cases:

 $\mathcal{M}(\rho) = R_{\mathcal{M}} * \rho$ quantum operation

 $\operatorname{Tr}[P^t \rho] = P * \rho$ POVM

$\rho \otimes \sigma = \rho * \sigma$

tensor product

$\mathrm{Tr}_{\mathsf{H}}[R] = R * I_{\mathsf{H}}$

partial trace



PRL 101 060401 (2008) Gircuits Architecture Optimization



in-out

The Choi operators of a fixed inputoutput comb structure make a convex set



PRL 101 060401 (2008) Circuits Architecture Optimization

in-out

- The Choi operators of a fixed inputoutput comb structure make a convex set
- Causality constraints correspond to a hyperplane section of the convex
- Group-covariance gives another linear constraint:

$[R, V_g] = 0 \implies R = \bigoplus R_j \otimes \mathbb{1}_{m_j}$







The mathematical formulation is reduced to a convex problem!



Realization theorem

Chiribella, D'Ariano, Perinotti, PRL **101** 060401 (2008) EL **83** 30004 (2008) Theorem: Every Choi operator on given input-output spaces and satisfying given causality conditions is realized by the comb of isometries



For realization of isometries see: Buscemi, D'Ariano, and Sacchi, PRA 68 042113 (2003)



causality constraints:









Using quantum memory delay the use of subcircuits by breaking the comb into subcombs + quantum memory




Using quantum memory delay the use of subcircuits by breaking the comb into subcombs + quantum memory









Application 1: discrimination and estimation of unitaries (optimal oracle-calling quantum algorithms)



Chiribella, D'Ariano, Perinotti, PRL 101 180501 (2008)

Optimal discrimination between two possible unitary operators $U_1 U_2$





Chiribella, D'Ariano, Perinotti, PRL 101 180501 (2008)

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Chiribella, D'Ariano, Perinotti, PRL 101 180501 (2008)

Optimal discrimination between two possible unitary operators $U_1 U_2$



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spread lemma: $\Delta(AB) \leq \Delta(A) + \Delta(B)$ A M Childs, J Preskill, and J Renes, JMO 47, 155-176 (2000).

 $\Delta[W(U \otimes I)W^{\dagger}(U \otimes I)] \leq \Delta(U^{\otimes 2})$

The spread of the tester is not larger than that of $U^{\otimes N}$ and U^N



Chiribella, D'Ariano, Perinotti, PRL 101 180501 (2008)

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The spread of the tester is not larger than that of $U^{\otimes N}$ and U^N . The parallel diposition is already optimal



There are memory channels that can be discriminated perfectly with a single use by a quantum tester, and not conventionally



Discrimination of memory channels

Chiribella, D'Ariano, Perinotti, (unpublished)

Tester Born rule:

$$\operatorname{Tr}[P_j R] = p_j, \qquad \sum_j P_j = \Xi$$

Perfect discrimination for:





What happens for more than two unitaries?

What happens for non-unitary channels?

Covariant estimation of unitaries

- Covariant unitary estimation problem (group of unitaries, Haar-distributed)
- Problem: find the optimal tester for estimating the group element

One can prove that the optimal tester is covariant:

 $T = \int_{G} \mathrm{d} \, g T_{g} \quad T_{h} = (U_{h}^{\otimes N} \otimes I) \Theta(U_{h}^{\dagger \otimes N} \otimes I) \quad \Rightarrow \quad [T, U_{h}^{\otimes N} \otimes I] = 0$ Then: $T^{\frac{1}{2}}(|U_{g}\rangle\rangle\langle\langle\langle U_{g}|)^{\otimes N}T^{\frac{1}{2}} = T^{\frac{1}{2}}(U_{g}^{\otimes N} \otimes I)|I\rangle\rangle\langle\langle I|(U^{\dagger \otimes N} \otimes I)T^{\frac{1}{2}} = (U_{g}^{\otimes N} \otimes I)T^{\frac{1}{2}}|I\rangle\rangle\langle\langle I|T^{\frac{1}{2}}(U^{\dagger \otimes N} \otimes I)$

Any covariant tester is equivalent to a parallel scheme



Oracle-calling quantum algorithm as optimal discrimination of unitaries





Oracle-calling quantum algorithm as optimal discrimination of unitaries



Hidden-subgroup algorithms (Deutsh-Jozsa, Simon, etc): parallel calls are optimal



Oracle-calling quantum algorithm as optimal discrimination of unitaries



Hidden-subgroup algorithms (Deutsh-Jozsa, Simon, etc): parallel calls are optimal

Algorithms with no hidden subgroup (e.g. Grover) need a comb [C. Zalka, PRA 60, 2746 (1999)]



Oracle-calling quantum algorithm as optimal discrimination of unitaries



Hidden-subgroup algorithms (Deutsh-Jozsa, Simon, etc): parallel calls are optimal

Algorithms with no hidden subgroup (e.g. Grover) need a comb [C. Zalka, PRA 60, 2746 (1999)]

Q-combs: systematic method to determine optimal oracle-calling algorithms



When do we need a tester (not just a parallel discrimination):



When do we need a tester (not just a parallel discrimination):

For discrimination of non-unitary channels

For discrimination within non-covariant sets

For discrimination of memory channels



Existence of optimal non parallel optimal discrimination schemes

Quantum information Quantum information PRL 101 180501 (2008)

Existence of optimal non parallel optimal discrimination schemes

The proper distance for memory channels must be defined in terms of optimal discriminating testers

$$D(\mathscr{C}^{(N)}, \mathscr{D}^{(N)}) := \max_{\Xi^{(N)}} \left\| \left(I \otimes \Xi^{(N)\frac{1}{2}} \right) \Delta \left(I \otimes \Xi^{(N)\frac{1}{2}} \right) \right\|_{1}$$
$$\Delta := C - D$$

Operational network distance PRL 101 180501 (2008)

Existence of optimal non parallel optimal discrimination schemes

The proper distance for memory channels must be defined in terms of optimal discriminating testers

$$D(\mathscr{C}^{(N)}, \mathscr{D}^{(N)}) := \max_{\Xi^{(N)}} \left\| \left(I \otimes \Xi^{(N)\frac{1}{2}} \right) \Delta \left(I \otimes \Xi^{(N)\frac{1}{2}} \right) \right\|_{1}$$
$$\Delta := C - D$$

CB-norm distance only accounts for parallel discrimination schemes



Application 2: quantum protocols (cryptography, game-theory)



Quantum protocols



Quantum combs describe the most general strategies in multi-party protocols and games

G. Gutoski and J. Watrous, Proc. STOC, 565-574, (2007)





Quantum bit commitment



Quantum combs is the most suitable mathematical formulation of Alice and Bob's strategies in a quantum bit commitment protocol



Sketch of impossibility proof

Chiribella, D'Ariano, Perinotti, Schlingemann, and Werner (in preparation)



Thm: a QBC concealing protocol cannot be binding

Proof: continuity of the comb-Stinespring versus the operational distance between strategies



Application 3: Quantum-algorithm learning



Problem: run an unknown unitary that is available today on a quantum state that will be available tomorrow



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Problem: run an unknown unitary that is available today on a quantum state that will be available tomorrow



Alice owns quantum circuit that performs a very valuable algorithm U that she wants to keep undisclosed.



Problem: run an unknown unitary that is available today on a quantum state that will be available tomorrow



Alice owns quantum circuit that performs a very valuable algorithm U that she wants to keep undisclosed.



Bob needs to run Alice's algorithm on an input state that will be available tomorrow, but he can borrow the circuit from Alice only today for just a limited number of uses N, and with the circuit sealed.



Problem: run an unknown unitary that is available today on a quantum state that will be available tomorrow





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im information theory group

The only thing that Bob can do today, with the circuit available, is to use it on a input state known to him.



Problem: run an unknown unitary that is available today on a quantum state that will be available tomorrow



a information theory group

The only thing that Bob can do today, with the circuit available, is to use it on a input state known to him.

After that the only thing that remains available to Bob for tomorrow is the output state, which Bob can store on a quantum memory.



Problem: run an unknown unitary that is available today on a quantum state that will be available tomorrow



- The only thing that Bob can do today, with the circuit available, is to use it on a input state known to him.
- After that the only thing that remains available to Bob for tomorrow is the output state, which Bob can store on a quantum memory.



Therefore, Bob needs a quantum device that is capable of retrieving from the output state, namely recovering U and then running it on a new unknown state.



Problem: run an unknown unitary that is available today on a quantum state that will be available tomorrow





Problem: run an unknown unitary that is available today on a quantum state that will be available tomorrow



- Exact storing of quantum states is possible (quantum memory is a technological problem)
- Perfect storing of undisclosed unitaries over a quantum state is impossible even in-principle (Nielsen-Chuang no-programming theorem)





Problem: run an unknown unitary that is available today on a quantum state that will be available tomorrow



retrieving

storing
Problem: a black box computes an unknown function y = f(x)We can evaluate f on a finite set of points x_1, \ldots, x_N getting outcomes y_1, \ldots, y_N

ntum information theory group



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tum information theory group

$$f - y_1 y_2 \dots y_N$$

Problem: a black box computes an unknown function y = f(x)We can evaluate f on a finite set of points x_1, \ldots, x_N getting outcomes y_1, \ldots, y_N

$$f - y_1 y_2 \dots y_N$$

Subsequently, we are asked to compute f on a new point x , without using the black box f(x) = ?

Problem: a black box computes an unknown function y = f(x)We can evaluate f on a finite set of points x_1, \ldots, x_N getting outcomes y_1, \ldots, y_N

$$f - y_1 y_2 \dots y_N$$

Subsequently, we are asked to compute f on a new point x, without using the black box f(x) = ?

In classical computer science, statistical learning provides a method to solve this problem

Classical networks for learning:

Comparing x with f(x) for N times is not the only possibility: this just corresponds to the parallel configuration



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$$f_1 - y_1$$

 \vdots \vdots
 $f_N - y_N$

Classical networks for learning:

Comparing x with f(x) for N times is not the only possibility: this just corresponds to the parallel configuration

$$f_1 - y_1$$

 \vdots \vdots
 $f_N - y_N$

To learn better, one could use a sequential network:

$$f = g_1$$
 $f = g_2$ $f = g_3$

where g_1, g_2, \ldots, g_N are known functions

- Unknown function f \longrightarrow unknown quantum channel \mathcal{E}
- Classical program —> quantum network

tum information theory group

- Input X \longrightarrow quantum state ρ_{in}
- Output Y \longrightarrow quantum state ρ_{out}



Ouantum algorithm learning Ouansical guess Quantum "guess"

Physical implementation of the quantum guess: retrieving channel \mathcal{R} It retrieves the unknown transformation from the output state ρ_{out} and performs it on a new state ρ

 $\rho_{out} \to \hat{\mathcal{E}}$



 $Y \to \hat{f}$



Quantum algorithm learning Quantum "guess"

Physical implementation of the quantum guess: retrieving channel \mathcal{R} It retrieves the unknown transformation from the output state ρ_{out} and performs it on a new state ρ

 $\rho_{out} \to \hat{\mathcal{E}}$



 $Y \to \hat{f}$



ρ – Ε

Target: implementing / the unknown channel with maximum fidelity

find the optimal quantum comb



Figure of merit: input-output fidelity

$$F(\mathcal{E}, \hat{\mathcal{E}}) = \int \mathrm{d}\varphi \ F(\mathcal{E}(\varphi), \hat{\mathcal{E}}(\varphi)) \quad F(\rho, \sigma) = \mathrm{Tr}\left[(\rho^{\frac{1}{2}} \sigma \rho^{\frac{1}{2}})^{\frac{1}{2}}\right]$$

Consider the case where the set of channels is a group of unitary transformations.



Assuming a uniform prior for the unknown unitaries, we have the average fidelity

$$F = \int \mathrm{d}U \ F(\mathcal{U}, \mathcal{C}_U)$$



Comb of the learning network: $L = R * C_N * \cdots * C_2 * C_1 * \rho_{in}$

Fidelity: $F = \frac{1}{d^2} \int dU \ \langle \langle U | \langle \langle U^* |^{\otimes N} | L | U \rangle \rangle | U^* \rangle \rangle^{\otimes N}$

We can always optimize over covariant combs:

$$[L, U \otimes V^* \otimes U^{* \otimes N} \otimes V^{\otimes N}] = 0 \qquad \forall U, V$$







Decomposing the unitaries as

$$U^{\otimes N} \otimes I_A = \bigoplus (U_J \otimes I_{m_J})$$

J

one can prove that the optimal input states have the form

$$|\psi\rangle = \bigoplus_{J} a_J \frac{|I_J\rangle}{\sqrt{d_J}} \qquad a_J \ge 0$$

where $|I_J
angle
angle \in \mathcal{H}_J^{\otimes 2}$ is a maximally entangled state

This is the same form of the optimal states for estimation of the unknown unitary U with N copies

G Chiribella, G M D'Ariano, and M F Sacchi, Phys. Rev. A 72, 043448 (2005).

Bisio, Chiribella, D'Ariano, Facchini, Perinotti (unpublished)

Theorem: for any group of unitaries, for an input state of the optimal form

$$|\psi\rangle = \bigoplus_{I} a_{J} \frac{|I_{J}\rangle}{\sqrt{d_{J}}} \qquad a_{J} \ge 0$$

the optimal retrieving channel to extract U from the states

$$(U^{\otimes N} \otimes I_A) |\psi\rangle = \bigoplus_J a_J \frac{|U_J\rangle}{\sqrt{d_J}} \qquad a_J \ge 0$$

is achieved by a "measure-and-prepare" scheme. (estimation of the unknown unitary U: for outcome \hat{U} , just perform the unitary \hat{U})

For the optimal POVM, see G Chiribella, G M D'Ariano, and M F Sacchi, Phys. Rev. A 72, 043448 (2005).























Optimal retrieving is "measure-and-prepare": no need for quantum memory. We can measure immediately after applying U, and store the outcome \hat{U} in a classical memory.

We can make as many copies as we want (a quantum memory is degraded at every access).



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- N non-identical input unitaries (and/or non-identical target unitaries)
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- N non-identical input unitaries (and/or non-identical target unitaries)
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not necessarily optimal for

- learning general channels
- learning unitaries that do not form a group
- learning with restrictions on the available input states (entanglement)



Application 4: Optimal cloning of unitaries

G. Chiribella, G. M. D'Ariano, P. Perinotti PRL **101** 180504 (2008)

$$F = \int \mathrm{d}U \, F(\mathscr{T}_U^{(N)}, \mathscr{T}_U^{\otimes N}) \qquad \text{(channel fidelity)}$$



Cloning of unitaries

Chiribella, D'Ariano, Perinotti, PRL 101 180504 (2008)





Application 5: Optimal quantum

tomography



Optimal tomographers



(d⁴ outcomes)

Informationally complete tester



Optimal tomographers

(d⁴ outcomes)

Informationally complete tester





Optimal tomographers



Informationally complete tester







circuit board tomographer



Optimal tomography

Use different in and out dimensions to unify: states, channels, and POVMs

arXiv: 0806.1172



A. Bisio, G. Chiribella, G. M. D'Ariano, S. Facchini, P. Perinotti PRL 102 010404 (2009)


Optimal tomography

- Prior distribution of channels corresponding to the depolarizing average channel
- Cost function = representation, (equally weighted orthonormal set of operators)
- Further selection:
 1) quantum operations,
 2) channels,
 3) unital channels

Use different in and out dimensions to unify: states, channels, and POVMs





arXiv: 0806.1172

A. Bisio, G. Chiribella, G. M. D'Ariano, S. Facchini, P. Perinotti PRL 102 010404 (2009)



Optimal tomography

arXiv: 0806.1172





Conclusions

PRL **101** 060401 (2008) PRL **101** 180501 (2008) PRL **101** 180504 (2008) PRL **102** 010404 (2009) EL **83** 30004 (2008)



Conclusions



- New Quantum Estimation Theory, with multiple copies, and optimization of the setup \rightarrow optimization of quantum circuits architecture, engineering high-precision operations
 - Quantum circuit board = quantum comb = supermap
 - Comb algebra (link-product)
 - Convex optimization method
- Section Applications:
 - Ş

Ş

- discrimination/estimation of unitaries and memory channels (optimal quantum oracle-calling algorithms)
- quantum protocols
- quantum-algorithm learning = storing undisclosable unitaries
- cloning undisclosable unitaries
- process tomography